

AN EFFICIENT ALGORITHM FOR LINE CLIPPING BY CONVEX POLYGON

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Abstract—A new line clipping algorithm against convex window based on a new approach for intersection detection is presented. Theoretical comparisons with Cyrus-Beck's algorithm are shown together with experimental results obtained by simulations. The main advantage of the presented algorithm is the substantial acceleration of the line clipping problem solution and that edges can be oriented clockwise or anti-clockwise.

1. INTRODUCTION

Many algorithms for clipping lines against rectangular, convex or non-convex windows have been published with many modifications derived from well-known Cohen-Sutherland, Liang-Barsky and Cyrus-Beck's algorithms [1-37]. Nevertheless, new developments or speed-up can be seen within every new paper.

The proposed algorithm is based on a simple rule: *Make tests first and then compute.*

A new criterion for testing whether the line intersects the given polygon or not is given, too.

The Cohen-Sutherland's algorithm (C-S) for clipping lines against a rectangular window is very well-known, not only because of its frequent usage, but for a very clever test whether the end-points are inside of the given rectangular area. The situation is a little bit more complicated when edges of the window are not collinear with axes.

1.1. Algorithm 1

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procedure Clip 2D Cyrus Beck (  $x_A, x_B$  );
begin { !! all vectors  $n_i$  precomputed !! algorithm shortened }
     $i := 1$ ;  $t_{min} := -\infty$ ;  $t_{max} := \infty$ ;  $s := x_B - x_A$ ;
    { for line segment  $t_{min} := 0.0$ ;  $t_{max} := 1.0$  }
    while  $i \leq N$  do {  $N$  is a number of edges }
        begin {  $n_i$  is a normal vector, e.g.  $n_i = [s_y, -s_x]^T$  }
            { and  $n_i$  must point out of the convex window }
             $\xi := s^T n_i$ ;  $s_i := x_i - x_A$ ;
            if  $\xi < > 0.0$  then
                begin  $t := s_i^T n_i / \xi$ ;
                    if  $\xi > 0.0$  then  $t_{max} := \min(t, t_{max})$ 
                        else  $t_{min} := \max(t, t_{min})$ 
                end
            else Special case solution;
             $i := i + 1$ 
        end;
    if  $t_{min} > t_{max}$  then EXIT; { !!  $\langle t_{min}, t_{max} \rangle = \emptyset$  }
    { recompute end-points if changed }
    if  $t_{max} < 1.0$  then  $x_B := x_A + s t_{max}$ ;
    if  $t_{min} > 0.0$  then  $x_A := x_A + s t_{min}$ ;
    SHOW LINE(  $x_A, x_B$  );
end;
    
```

The known algorithms for clipping lines against general convex window do not make tests similar to C-S algorithm. The main reason seems to be the com-

putational cost of such a test for convex window. If clipping algorithm is to be effective it is necessary to distinguish cases where lines pass through a given window from those where lines do not intersect the window. Cyrus-Beck's (C-B) algorithm solves this problem by direct computation of points of intersections, see Algorithm 1.

2. PROPOSED ALGORITHM

It is obvious that the line p intersects the edge x_0x_1 of the given polygon *if and only if* the vector s lies between two vectors s_0 and s_1 , see Fig. 1.

Let us define ξ and η as the z -coordinate of the cross products as follows

$$\xi = [sxs_i]_z \quad \eta = [sxs_{i+1}]_z \quad i = 0, \dots, N-1$$

where N is a number of edges.

It is possible to show that the line p given by the end-point x_A and by the vector s intersects the edge $x_i x_{i+1}$ *if and only if* the expressions

$$(\xi > 0) \text{ xor } (\eta < 0), \text{ resp. } \xi * \eta < 0$$

have value *true*, see Table 1 and Fig. 2. The above mentioned expressions are *independent* on polygon *orientation*. It means that edges can be clockwise or anti-clockwise oriented.

Now, it is possible to specify a new algorithm that is based on the presumption that a line can intersect a convex polygon only in two points, see Algorithm 2, if special cases are not considered, see Fig. 2. Special cases should be handled carefully, see [38].

In comparison with Algorithm 1 the proposed algorithm *only detects* whether the given line intersects the given polygon and then computes intersection points, while C-B algorithm does *compute all possible* intersection points with all edges.

2.1. Algorithm 2

```

procedure Clip 2D (  $x_A, x_B$  );
begin { !!!! all vectors  $\hat{s}_i$  are precomputed !!!! }
   $k := 0; i := N - 1; j := 0;$ 
   $s := x_B - x_A; \xi := [s \times s_i]_z;$ 
  while (  $j < N$  ) and (  $k < 2$  ) do
    begin  $\eta := [s \times s_j]_z;$ 
      if (  $\xi * \eta < 0.0$  ) then { intersection exists }
        begin
           $index_k := i; \{ \text{save edge index having intersection} \}$ 
           $k := k + 1; \xi := \eta;$ 
        end
      else if (  $\xi * \eta = 0.0$  ) then Special cases {  $\xi=0$  or  $\eta=0$  }
        { cases i - m }
        else Intersection does not exist;
        { cases e - h }

       $i := j; j := j + 1$ 
    end;
    if  $k = 0$  then { intersection does not exist } EXIT;
    {  $k = 2$  intersections exist - edges saved in  $index_k$  }
     $t_{min} := -\infty; t_{max} := \infty; \{ \text{for line segments } t_{min} := 0; t_{max} := 1; \}$ 
     $i := index_1; t_1 := \det[ x_i - x_A \mid -\hat{s}_i ] / \det[ s \mid -\hat{s}_i ];$ 
    { for effective implementation use identity  $x_i - x_A = s_i$  }
     $i := index_2; t_2 := \det[ x_i - x_A \mid -\hat{s}_i ] / \det[ s \mid -\hat{s}_i ];$ 
    { recompute end-points if changed }
    if  $t_1 < t_2$  then { swap  $t_1 \leftrightarrow t_2$  values ? }
    begin if  $t_2 < t_{max}$  then  $x_B := x_A + s t_2;$ 
    if  $t_1 > t_{min}$  then  $x_A := x_A + s t_1$ 
    end {  $t_2 < t_{max}; t_1 < t_{min}$  can be left out for lines }
    else begin if  $t_1 > t_{min}$  then  $x_B := x_A + s t_1;$ 
    if  $t_2 < t_{max}$  then  $x_A := x_A + s t_2$ 
    end;
    SHOW LINE(  $x_A, x_B$  );
  end

```

3. THEORETICAL ANALYSIS

Before making any comparisons it is necessary to point out that time needed for each operation ($:=, <, \pm, *, /$) does differ from computer to computer, see Table 2.

Let us denote N number of edges of the given convex polygon. Then it is possible to express the complexity of the C-B algorithm as (the worst case)

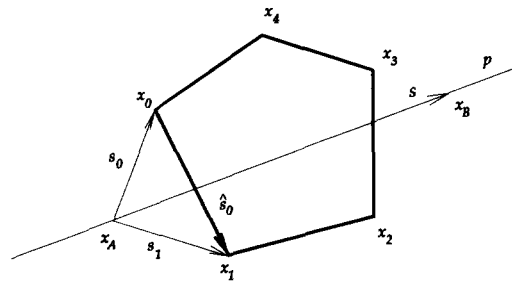


Fig. 1. Line segment clipping by convex polygon.

$$(8, 3, 6, 4, 0) + (5, 3, 7, 4, 1) * N$$

and time of computation, according to Table 2 (PC 486 used), can be estimated as

$$T_{C-B} = 590 + 621 * N$$

while for proposed algorithm the complexity is given as

$$(15, 3, 11, 14, 2) + (3, 1, 3, 3, 0) * N$$

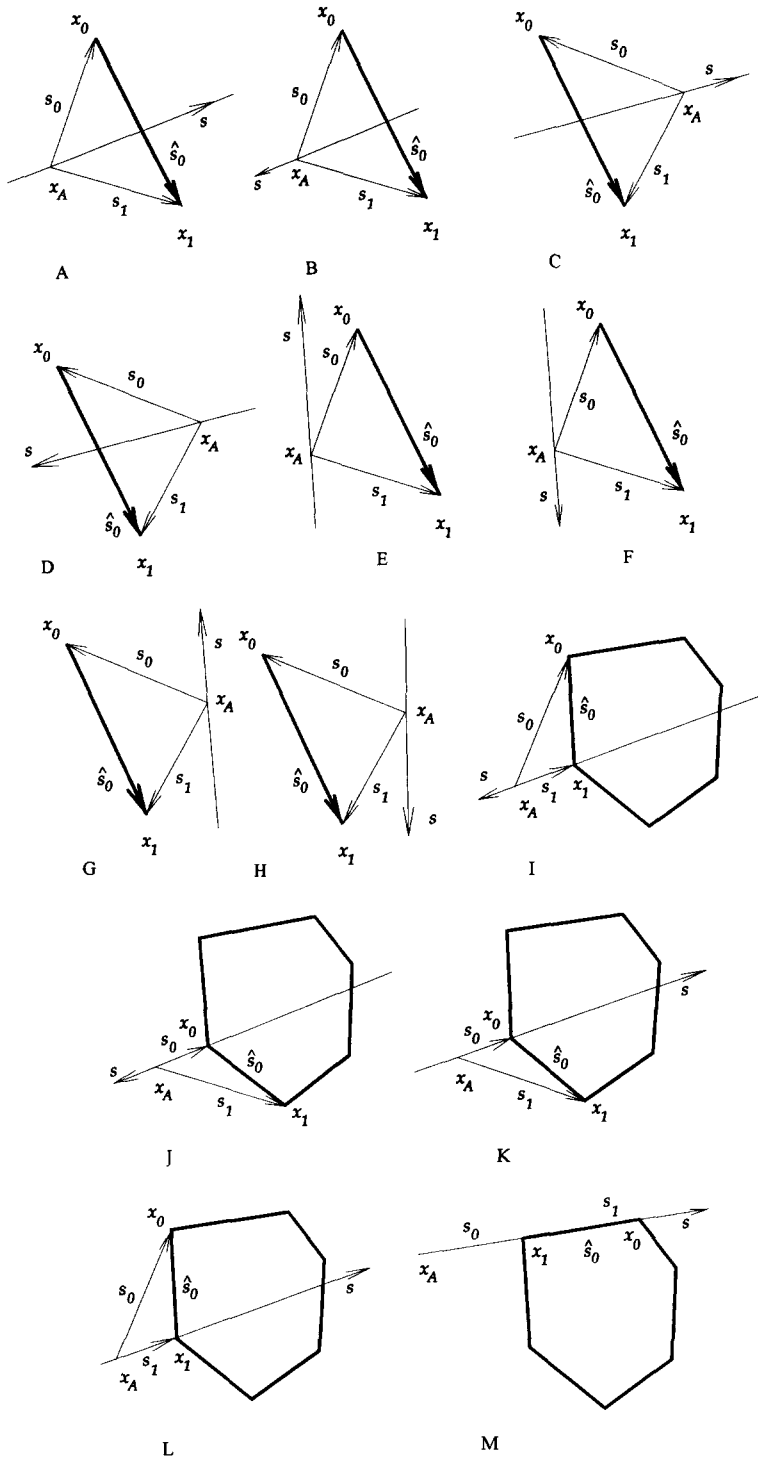


Fig. 2. Possible cases for line segment clipping.

and time of computation T can be estimated

$$T = 1329 + 257 * N$$

It means that the new algorithm should be better than C-B algorithm for $N \geq 2.1$. Let us consider a coefficient ν as a speed up coefficient, which is defined as

$$\nu = \frac{T_{C-B}}{T}$$

then some theoretical results of speed up are given in Table 3.

From the theoretical point of view the proposed algorithm should be for $N = 4$, approximately 30% faster than C-S algorithm. It can be seen that the efficiency increases dramatically with N , especially for $4 \leq N \leq 20$. For $N > 20$ the algorithm needs *only half* of the time needed by C-B algorithm. For $N > 50$ the factor ν is nearly constant. Table 3 shows expected results for

Table 1. Possible cases.

ξ	ν	Does the line intersect edge?
>0	>0	no—cases <i>g, h</i>
>0	=0	yes—special case <i>l</i>
>0	<0	yes—cases <i>a, c</i>
=0	>0	yes—special case <i>j</i>
=0	=0	no—special case <i>m</i>
=0	<0	yes—special case <i>k</i>
<0	>0	yes—cases <i>b, d</i>
<0	=0	yes—special case <i>i</i>
<0	<0	no—cases <i>e, f</i>

Table 2. Times for basic arithmetic operations.

	PC 386DX/387 25 MHz				
	:=	<	±	*	/
Int	14	42	10	40	58
Float	204	260	80	82	154

	PC 486 33 MHz 64KB cache—selected for evaluation				
	:=	<	±	*	/
Int	5	9	3	26	44
Float	33	50	16	20	114

Time is in 1/10 sec. for 5 000 000 operations.

Table 3. Theoretical results.

Expected results for PC 386DX/387							
	Intersections exist			Intersections do not exist			
	N	4	10	$\Rightarrow \infty$	4	10	$\Rightarrow \infty$
ν	1.26	1.60	2.09	1.75	1.87	1.97	

Expected results for PC 486							
	Intersections exist			Intersections do not exist			
	N	4	10	$\Rightarrow \infty$	4	10	$\Rightarrow \infty$
ν	1.30	1.74	2.42	2.01	2.16	2.29	

Table 4. Experimental results.

Experimental results for PC 386DX/387								
	5	8	10	15	20	30	70	190
	ν	1.41	1.54	1.73	1.81	1.90	2.08	2.27

Experimental results for PC 486 33 MHz								
	5	8	10	15	20	30	70	190
	ν	1.37	1.59	1.71	1.88	2.25	3.32	2.58

PC 386/387 DX, too. It can be seen that the efficiency of the proposed method should be higher for faster machines. It is necessary to point out that many simplifications have been made for theoretical analysis because time of operations for integers and time needed for instructions was not considered.

4. EXPERIMENTAL RESULTS

The C-B and proposed algorithms have been tested on data sets of end-points that have been randomly and uniformly generated over a space inside of a circle in order to eliminate an influence of rotation. Convex polygons were generated as regular N -sided convex polygons inscribed into a smaller circle. The results in Table 4 have been obtained if 80% of the lines intersect the window.

The coefficients ν are in a good correlation with theoretical results because the pessimistic estimations have been taken for the theoretical analysis. It is necessary to point out that different tests might be used instead of shown cross product, e.g., test for detection on which side of the given line p the point x_i lies.

5. CONCLUSION

The new efficient algorithm for clipping lines against convex window has been developed. Edges of the given polygon can be arbitrarily oriented. The theoretical analysis showed the advances over well-known Cyrus-Beck's algorithm. Experimental results support the theoretical analysis in spite of the fact that many factors have been omitted, like computational cost with integers, etc.

All tests were implemented in C++ on a PC 386/387 25 MHz and PC 486 33 MHz. It can be expected that for workstations the efficiency ν will be higher than for PC 486.

There is a hope that the proposed algorithm can be modified for clipping lines against non-convex windows, similarly to [31].

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